



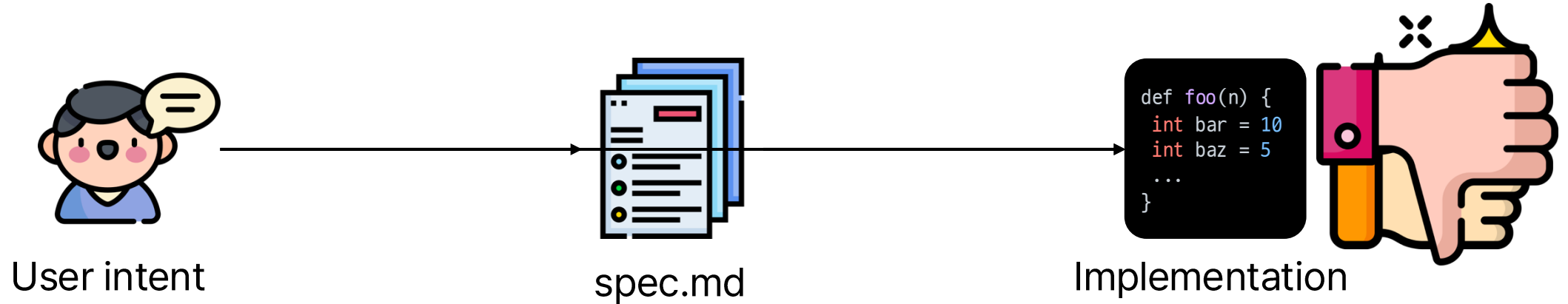
Expecto:

Extracting Formal Specifications from
Natural Language Description for Trustworthy Oracles

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Spec-Driven Development is Emerging



Many of AI coding tools support spec-driven development



Kiro (Amazon)



SpecKit (GitHub)

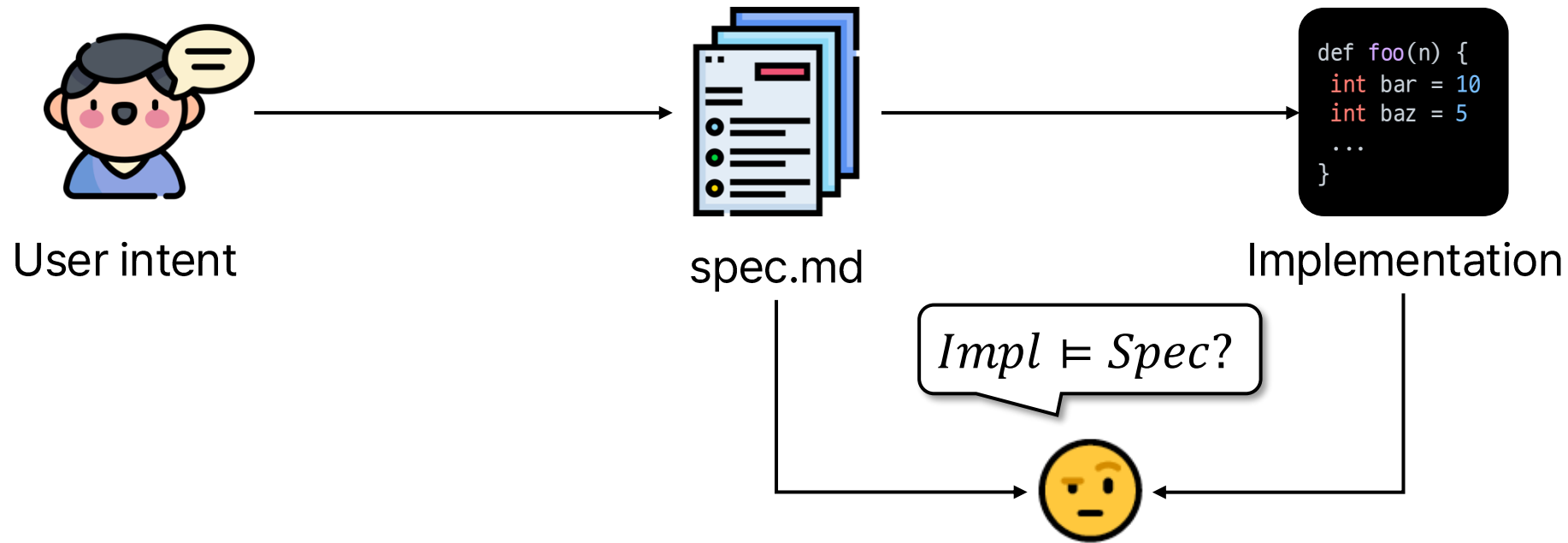


Codex (OpenAI)



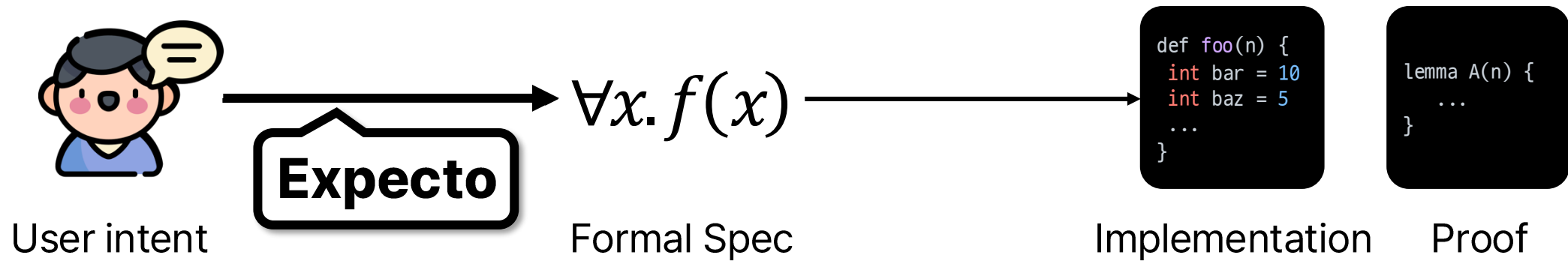
Claude Code (Anthropic)

Limitation of Spec-Driven Development



- How to check implementation satisfies natural language intent?
- 🤔 **Manual Inspection!** or simply asking an LLM...

Use Formal Verification! But, ...

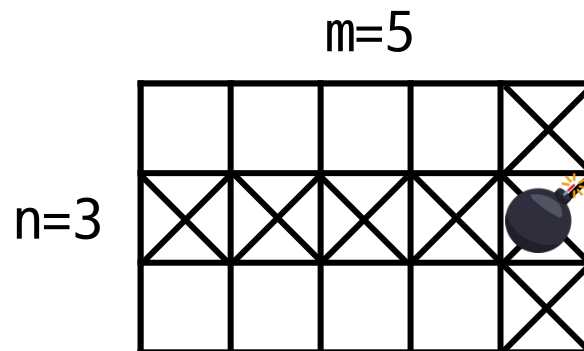


- Problem: writing spec in formal language is challenging
- **Solution: convert natural language intent into formal spec!**

Running Example: A Bomb (Codeforces 699B)

Natural Language Intent

- A **field** is an $n \times m$ grid consisting of walls (“X”) and empty cells (“ ”).
- If you place a bomb at (x, y) , all walls in row x and column y are removed.
- Determine whether it is possible to remove all walls with a single bomb, and if so, return the values of x and y .



Possible, $x=2$, $y=5$

Problem of Monolithic Spec Generation

- Existing methods generate formal spec at once (Endres et al.)

```
assert possible == any(  
    all((cell[i][x - 1] == "*" or cell[y - 1][j] == "*")  
        if cell[i][j] == "*" else True  
        for i in range(n)  
        for j in range(m))  
    for y in range(1, n + 1)  
    for x in range(1, m + 1)  
) and (  
    not possible or  
    all((cell[i][x - 1] == "*" or cell[y - 1][j] == "*")  
        if cell[i][j] == "*" else True  
        for i in range(n)  
        for j in range(m))  
    for y, x in [(y, x)]  
)
```

Wrong condition

Duplicated Codes

Unnecessary Loop

Over-simplified condition

```
assert (not possible)  
    or (1 <= y <= n and 1 <= x <= m)
```

Expecto Generates Correct & Modular Spec

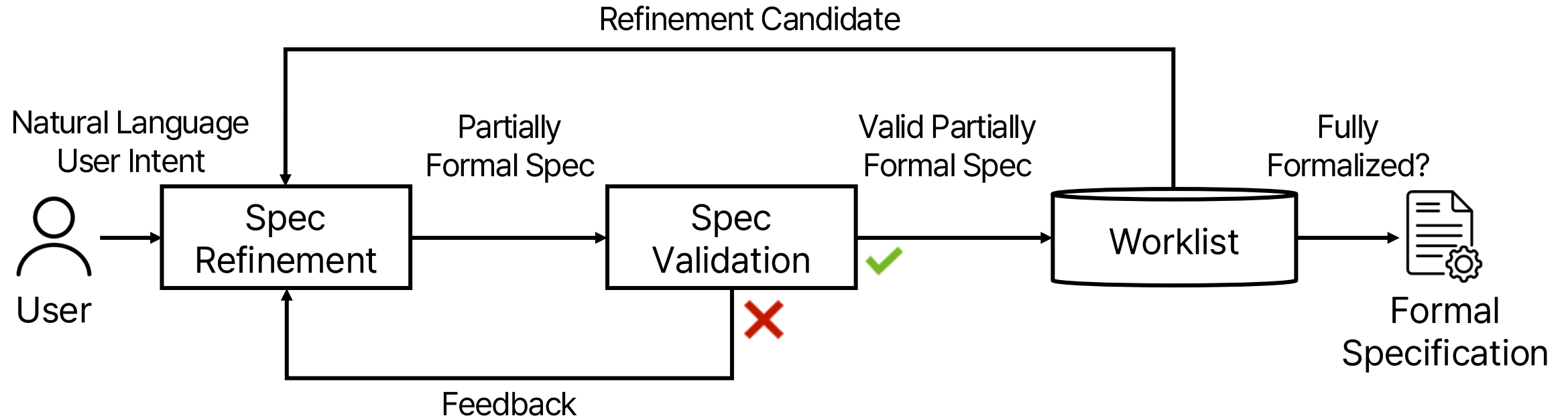
- A field is an $n \times m$ grid consisting of walls ("X") and empty cells (" ").
- If you place a bomb at (x, y) , all walls in row x and column y are removed.
- Determine whether it is possible to remove all walls with a single bomb, and if so, return the values of x and y .

At most **2.81x** more correct specs

Easy to understand specs

- $\text{Spec}(n, m, \text{field}, \text{possible}, x, y) =$
 $(\text{possible} \rightarrow (\text{ValidPos}(x, y, n, m) \wedge \text{WipesAll}(x, y, n, m, \text{field})) \wedge$
 $(\neg \text{possible} \rightarrow (\forall r, c. \text{ValidPos}(r, c, n, m) \rightarrow \neg \text{WipesAll}(r, c, n, m, \text{field})))$
- $\text{ValidPos}(r, c, n, m) = (1 \leq r \leq n) \wedge (1 \leq c \leq m)$
- $\text{WipesAll}(r, c, n, m, \text{field}) = \forall i, j. \text{IsWall}(i, j, n, m, \text{field}) \rightarrow (i+1=r \vee j+1=c)$
- $\text{IsWall}(r, c, n, m, \text{field}) = (0 \leq r < n) \wedge (0 \leq c < m) \wedge (\text{field}[r][c] = "X")$

Overview



- **Top-down spec synthesis** generates spec from outline to detail
- **Continuous spec validation** checks partial specs during synthesis

Top-Down Spec Synthesis

- D_i : informal definitions
- D_f : formal definitions



D_i Spec($n, m, \text{field}, \text{possible}, x, y$): "A field is an $n \times m$ grid..."

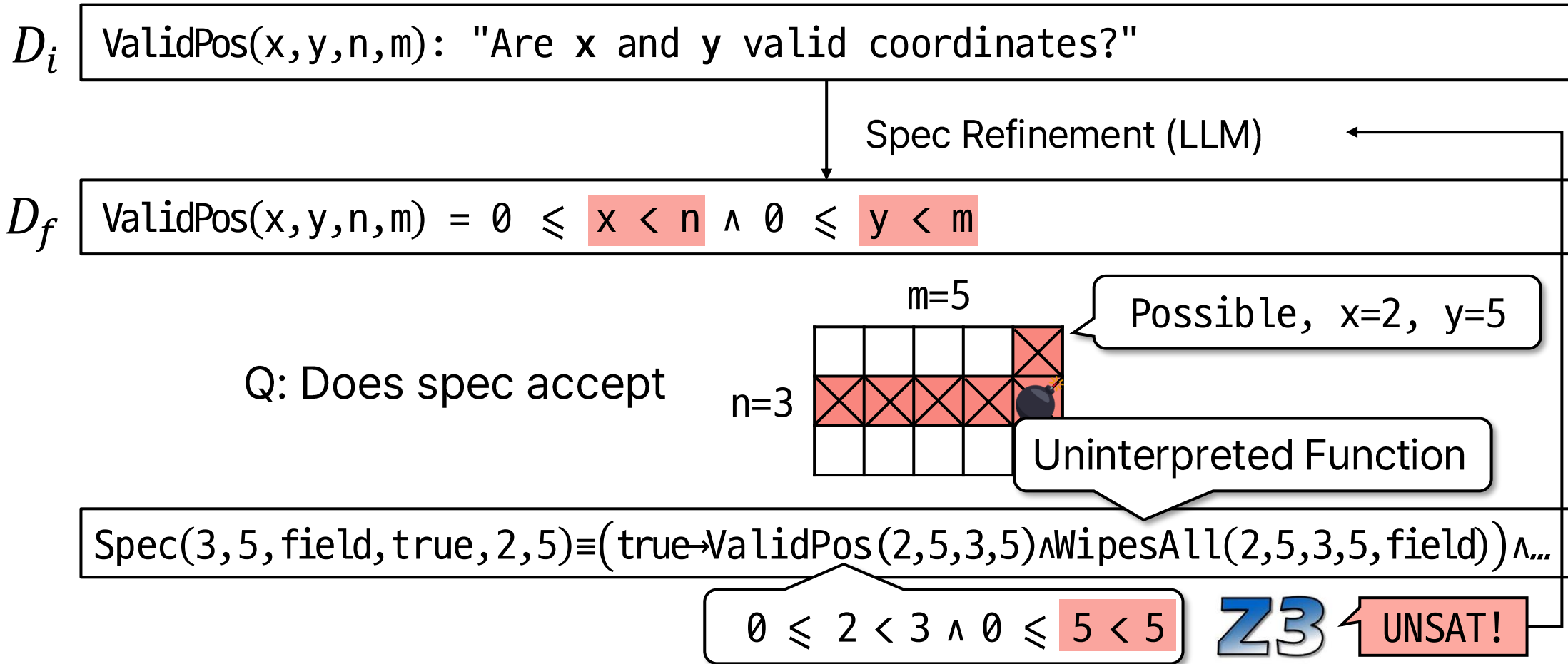


Spec Refinement (LLM)

D_f Spec($n, m, \text{field}, \text{possible}, x, y$) =
(possible $\rightarrow$ (**ValidPos**(x, y, n, m) $\wedge$ **WipesAll**(x, y, n, m, field)) $\wedge$
($\neg$ possible $\rightarrow$ ($\forall r, c. \text{ValidPos}(r, c, n, m) \rightarrow \neg \text{WipesAll}(r, c, n, m, \text{field})$))

- D_i
- ValidPos(x, y, n, m): "Are x and y valid coordinates?"
 - WipesAll(r, c, n, m, field): "For r and c , will all walls be removed?"

Continuous Spec Validation



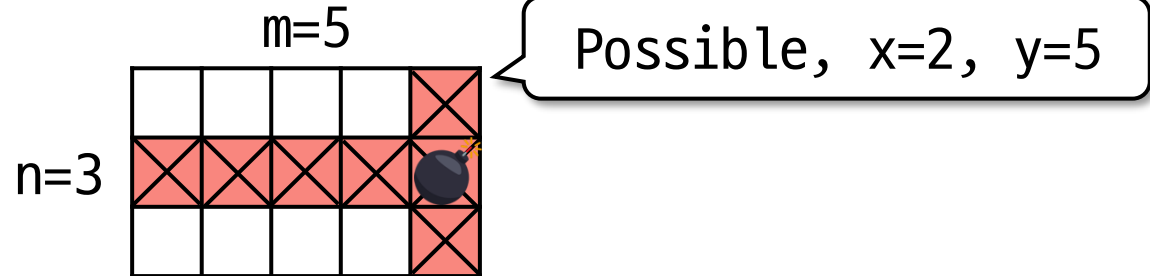
Continuous Spec Validation

D_i ValidPos(x,y,n,m): "Are x and y valid coordinates?"

Spec Refinement (LLM)

D_f ValidPos(x,y,n,m) = $0 < x \leq n \wedge 0 < y \leq m$

Q: Does spec accept



$\text{Spec}(3,5, \text{field}, \text{true}, 2,5) \equiv (\text{true} \rightarrow \text{ValidPos}(2,5,3,5) \wedge \text{WipesAll}(2,5,3,5, \text{field})) \wedge \dots$

$0 < 2 \leq 3 \wedge 0 < 5 \leq 5$

Z3

SAT!

Types of Validators

1. Syntax and type validators
2. Logical consistency validator
 - Is spec neither **tautology** nor **contradictory**?
3. Input-output example consistency validator
 - Does spec **accept correct** user-provided input-output examples?
 - Does spec **reject incorrect** user-provided input-output examples?

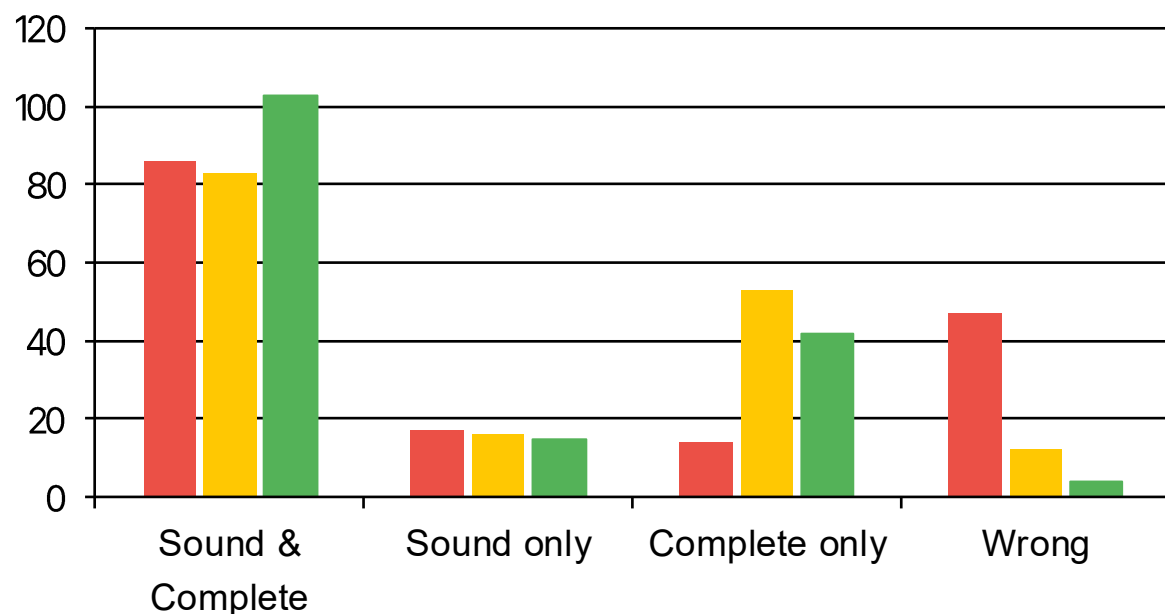
Experiment Setup

- Benchmark
 - Algorithmic problem: HumanEval+ (Liu et al.), APPS (Hendrycks et al.)
 - Real-world software: Defects4J (Just et al.)
- Baseline
 - NL2Postcond: Monolithic spec generation using LLM (Endres et al.)
- Model
 - GPT-4.1-mini (temperature=0.2)
- SMT Solver
 - Z3

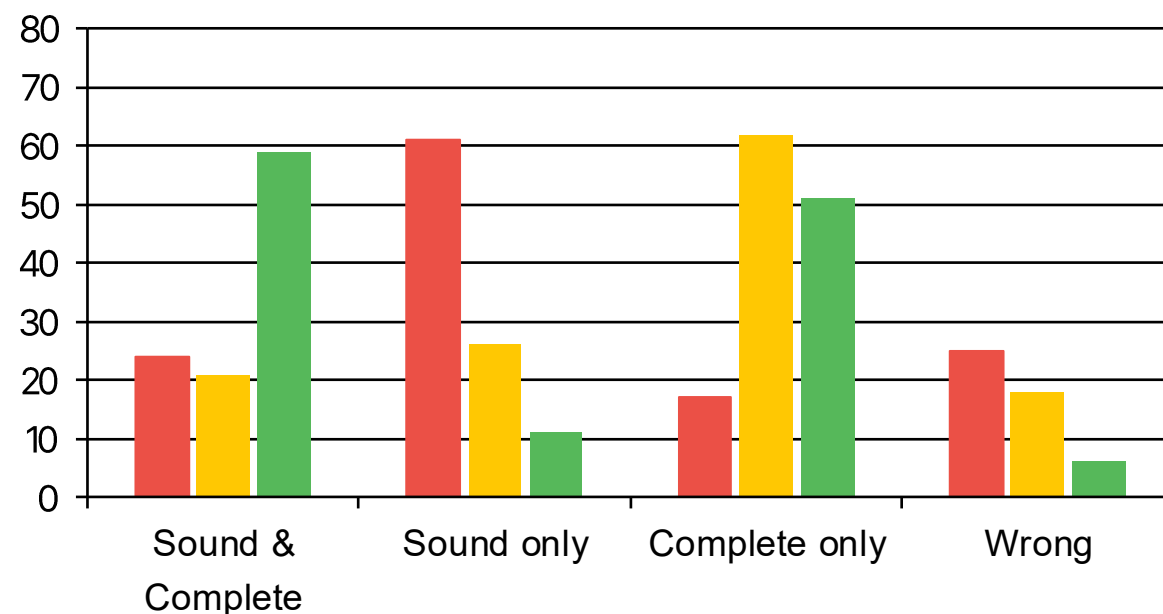
Can Expecto Generate Correct Specs?

NL2P NL2Post Full Expecto capturing a single constraint

Base Simple Expecto



Base Simple Expecto

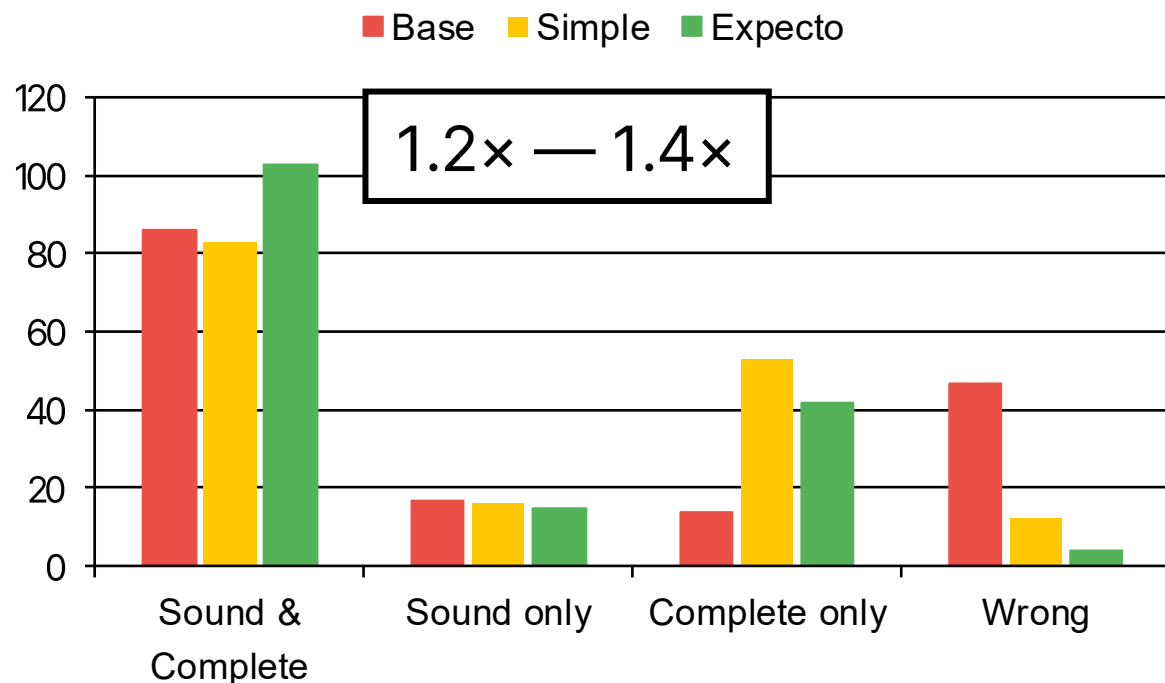


Soundness: Rejects incorrect input-output pairs

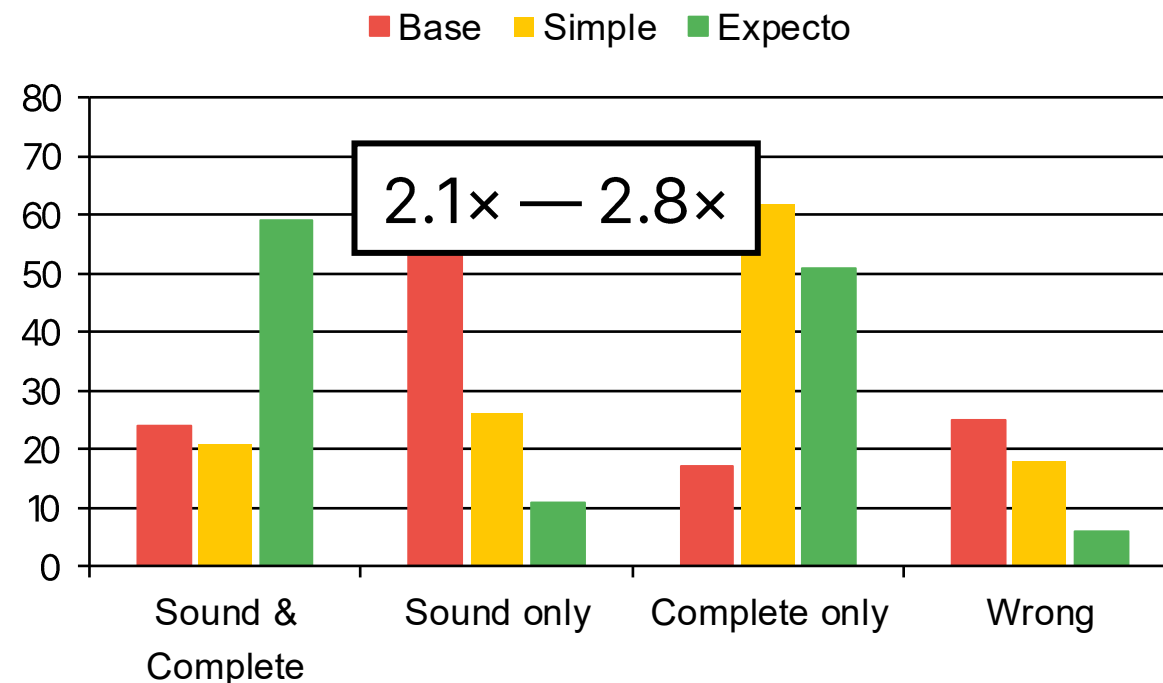
Completeness: Accepts correct input-output pairs

Can Expecto Generate Correct Specs?

HumanEval+



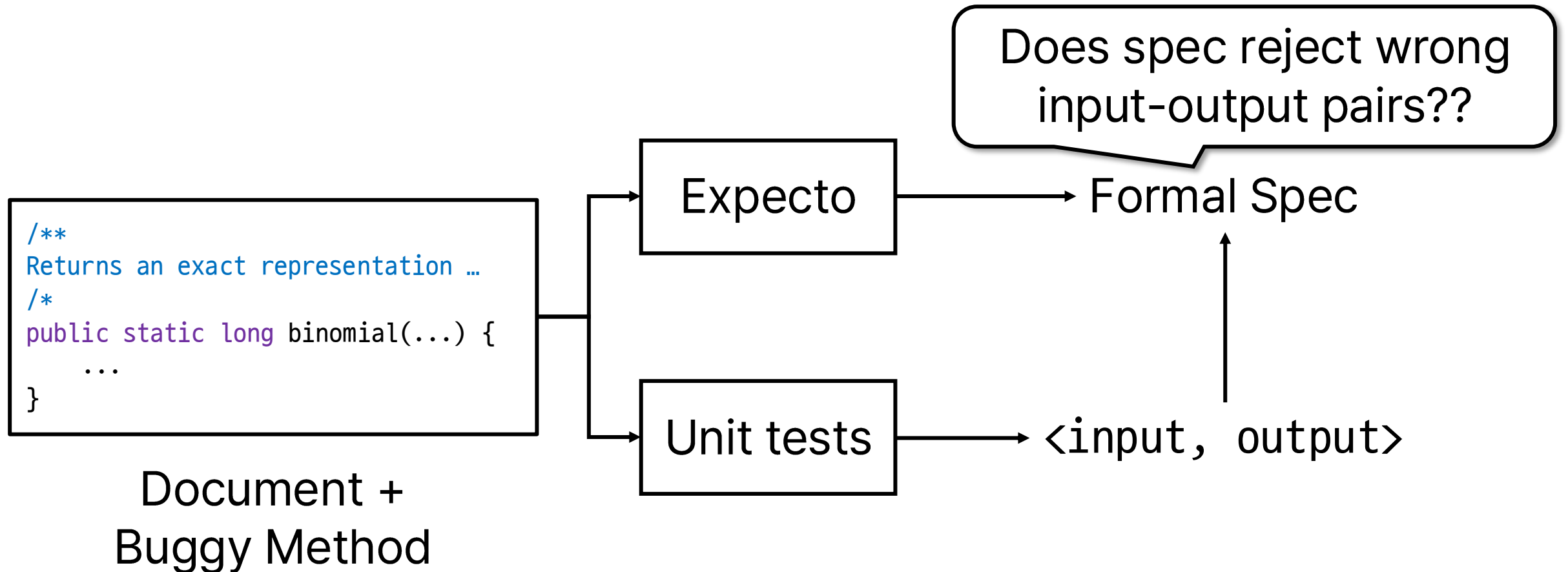
APPS



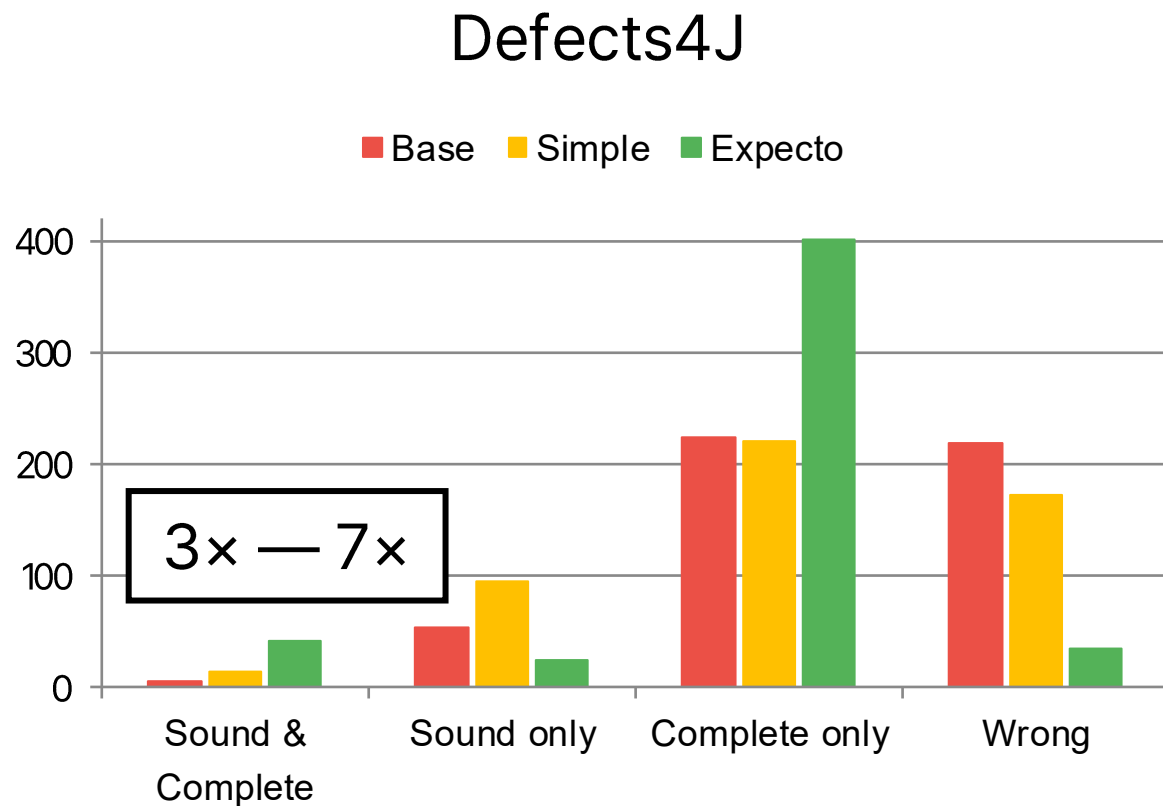
Expecto generates the most correct specs!

Can Expecto Detect Real-World Bugs?

- Detect bugs in Java methods using formal spec from document



Can Expecto Detect Real-World Bugs?



Expecto can detect bugs using generated specs

Summary

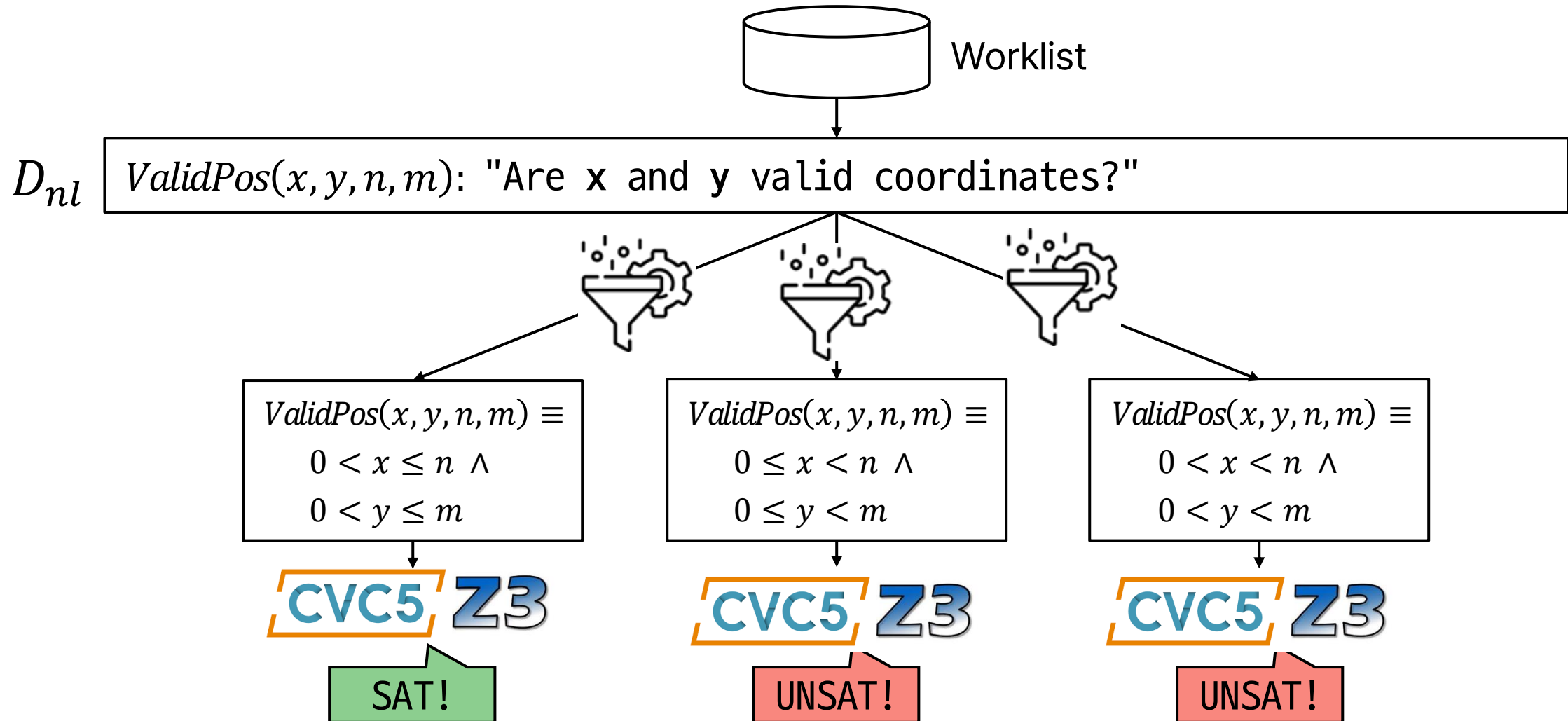


- **Problem: formal verification is promising but needs formal spec!**
 - Writing formal spec is challenging
 - Existing monolithic methods are limited
- **Solution: Expecto**
 - Top-down spec synthesis
 - Continuous spec validation
- **Results**
 - Generates more sound & complete specs on algorithmic (at most 2.8×) and real-world software benchmarks (at most 7×)

Homepage

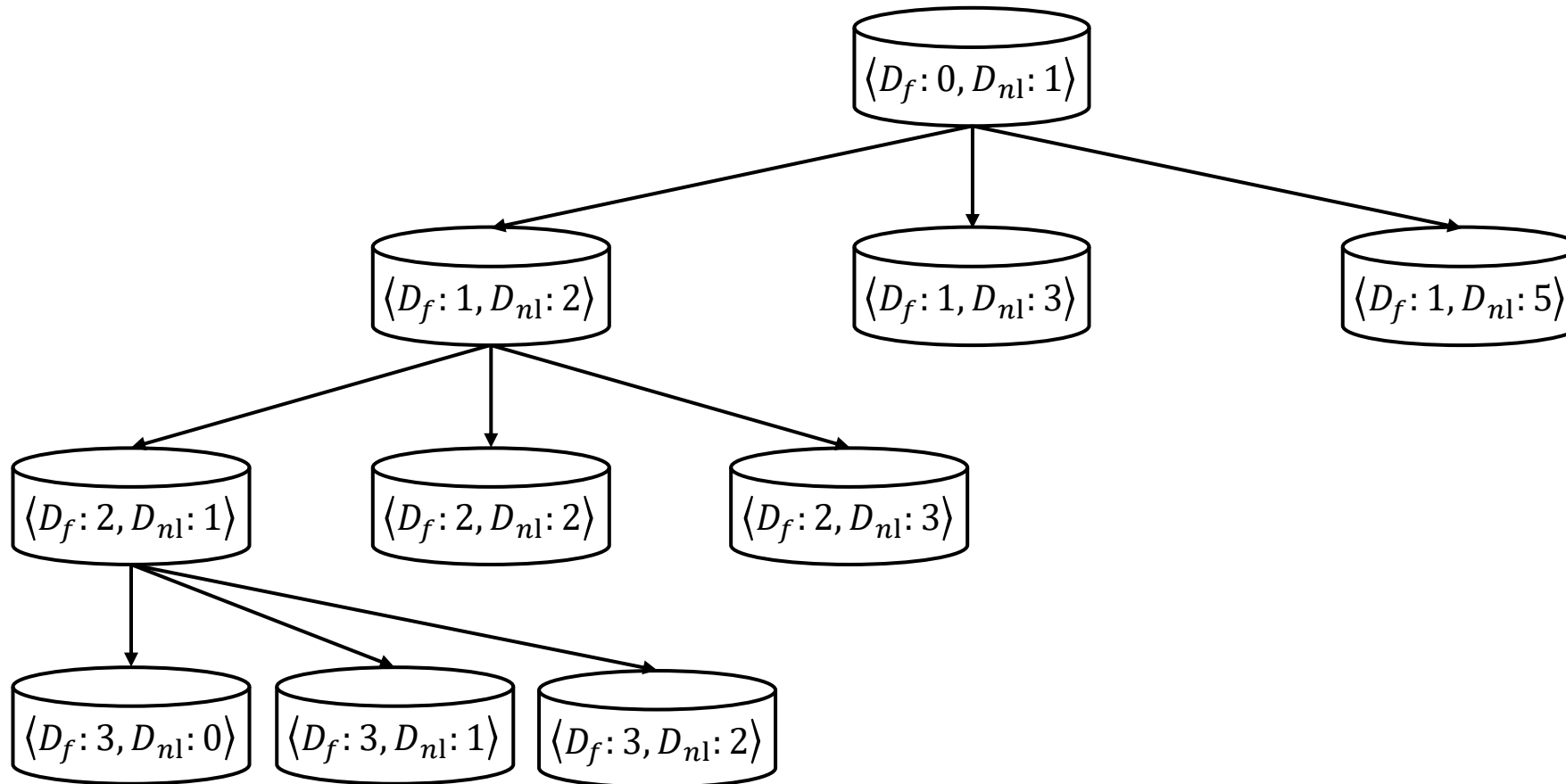


Tree Search Explores Diverse Paths



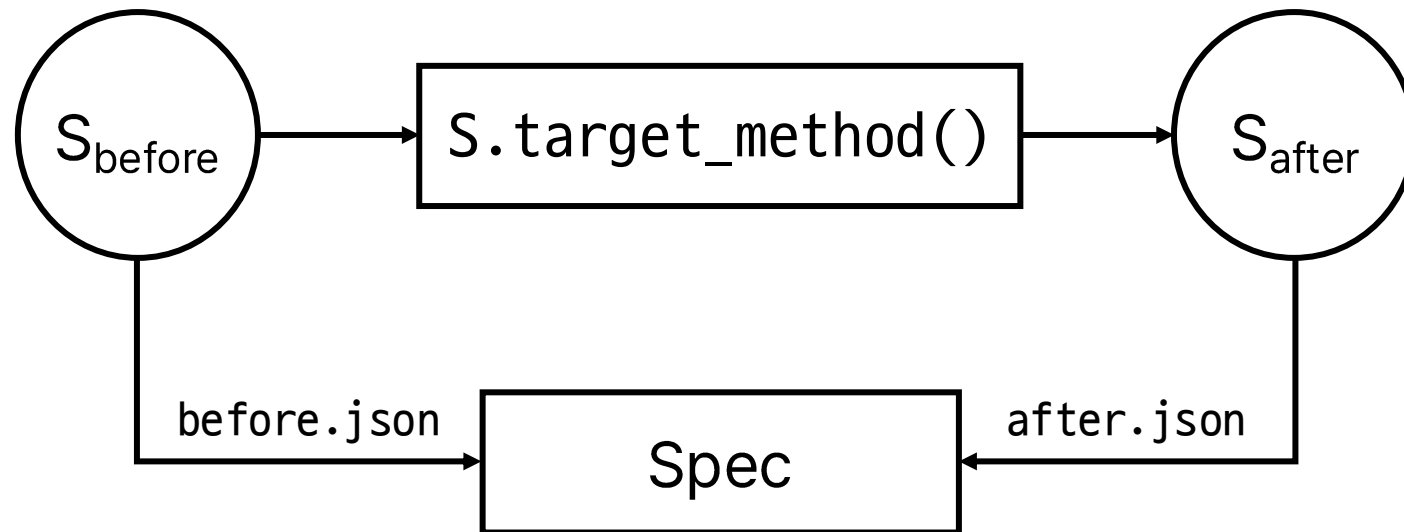
Tree Search prioritization strategy

- Larger D_f , smaller D_{nl} is better!



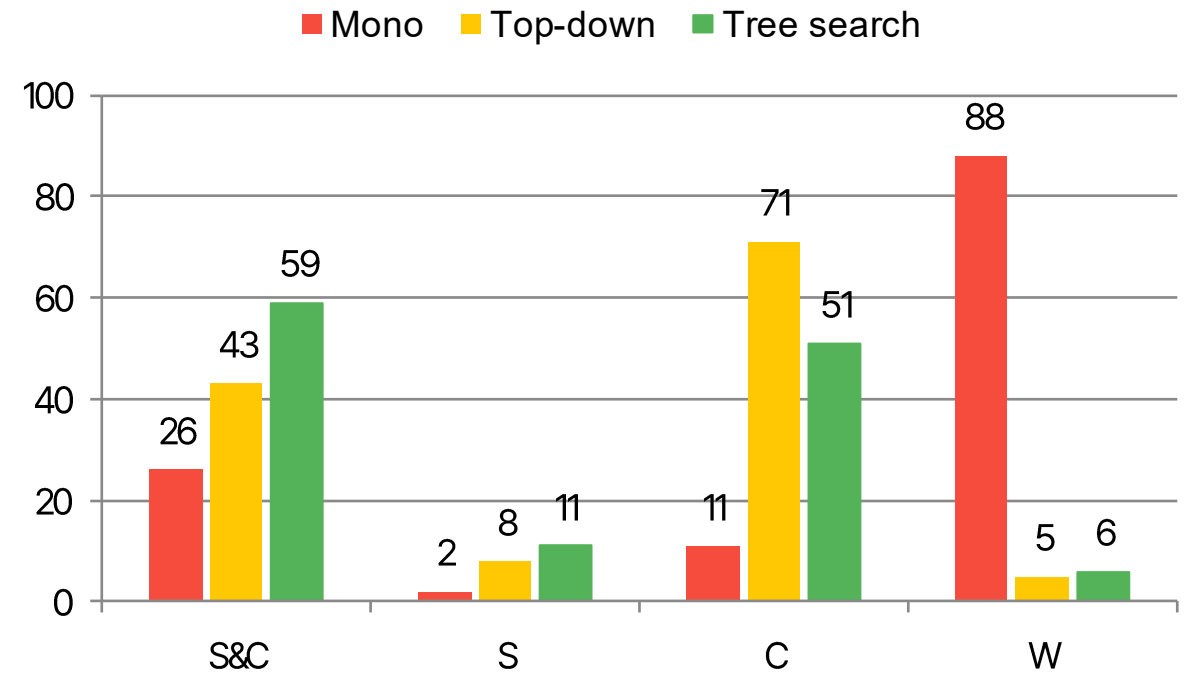
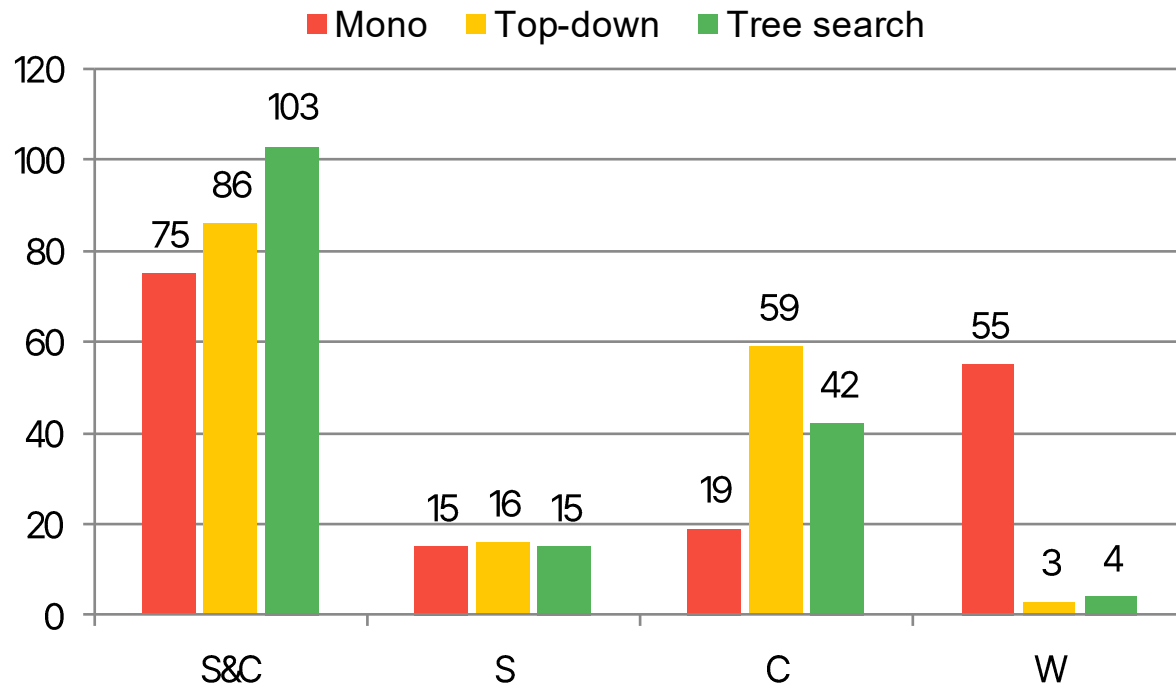
Modeling Side Effects

- Detect bugs in Java methods using formal spec from document
- But Java method is not pure!
- Solution: serialize states before and after execution



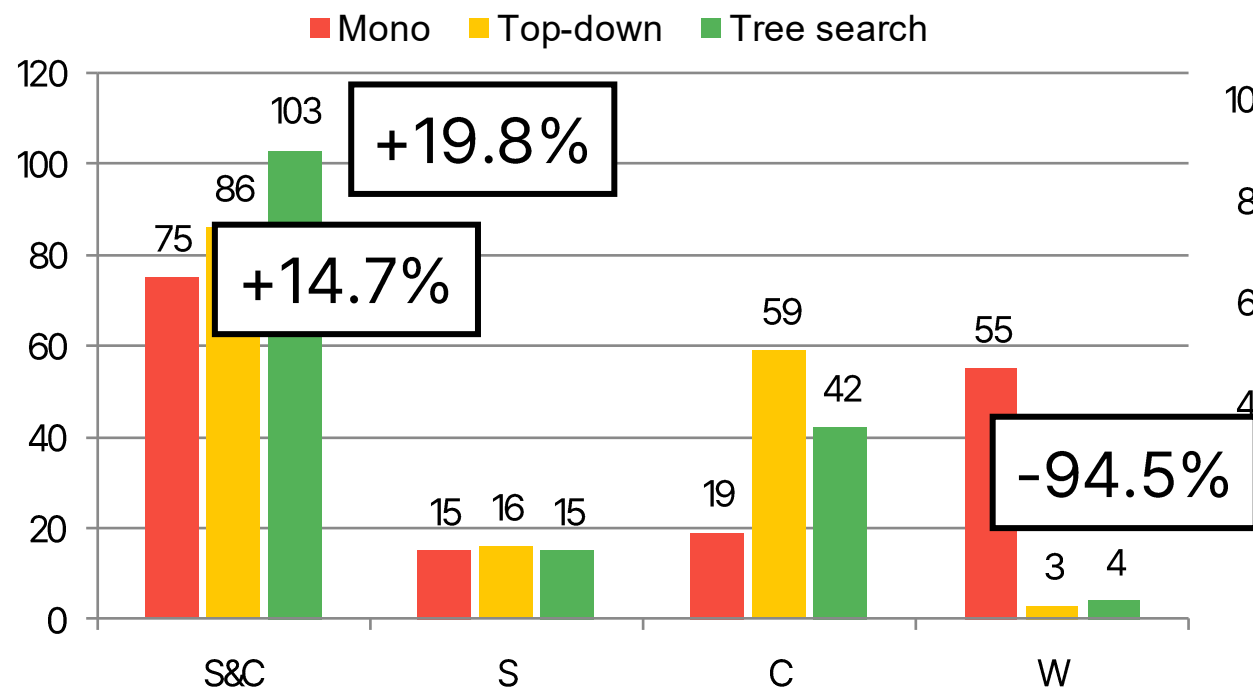
Impact of Search Algorithm

Ge Gener Generate spec with top-down + tree search ps

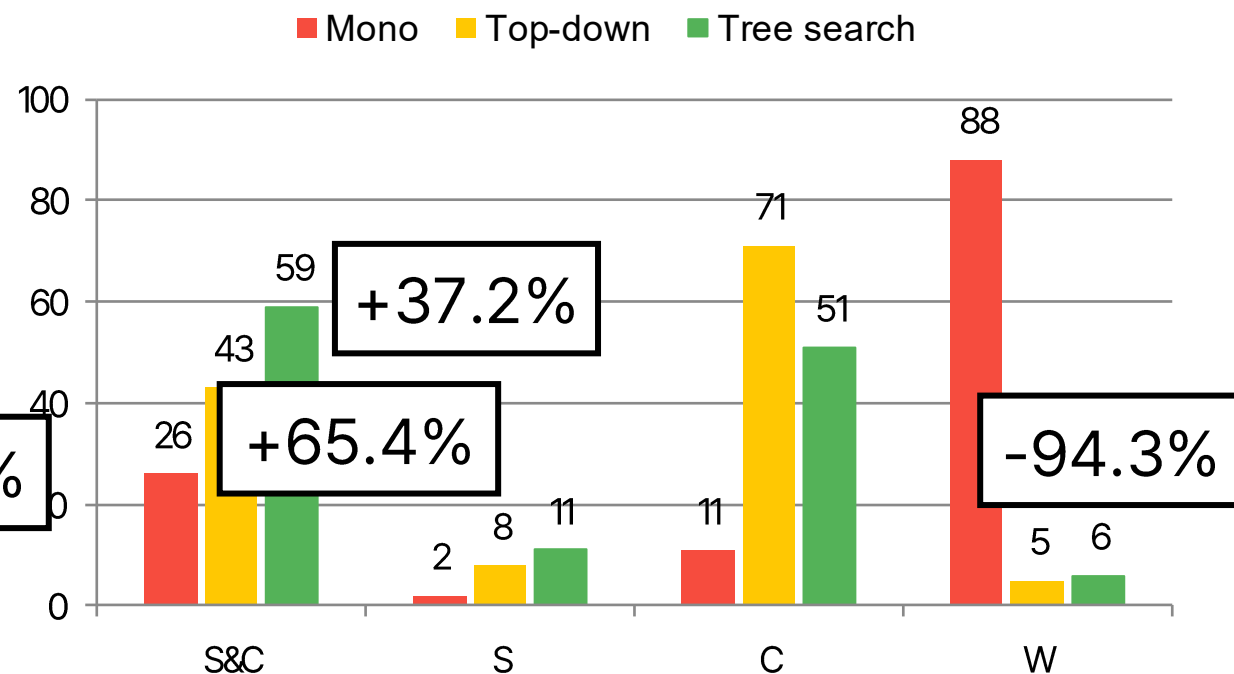


Impact of Search Algorithm

HumanEval+



APPS



1. Top-down spec synthesis reduces cognitive burden on LLM
2. Tree search explores diverse paths and it leads to more S&C

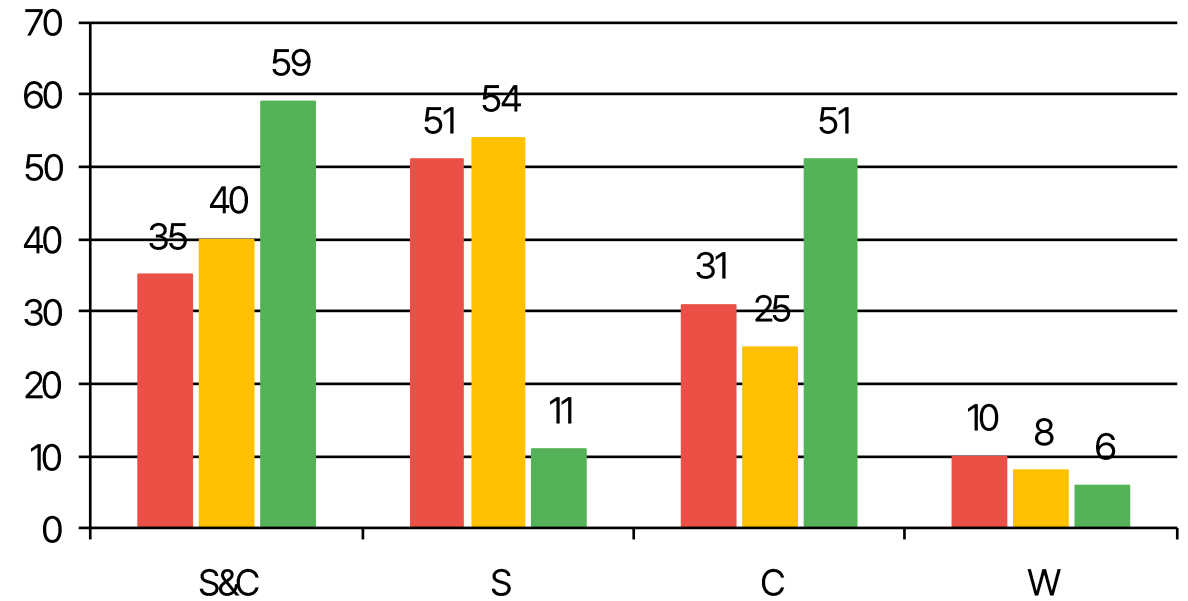
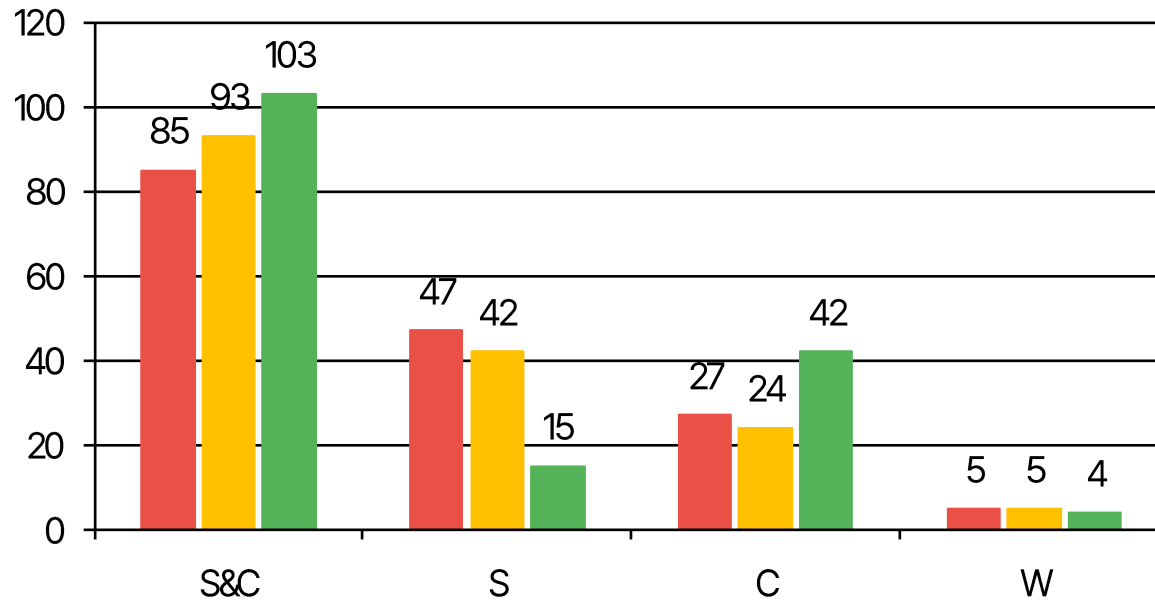
Impact of Checkers

S Van LC + test case consistency checker

APPS

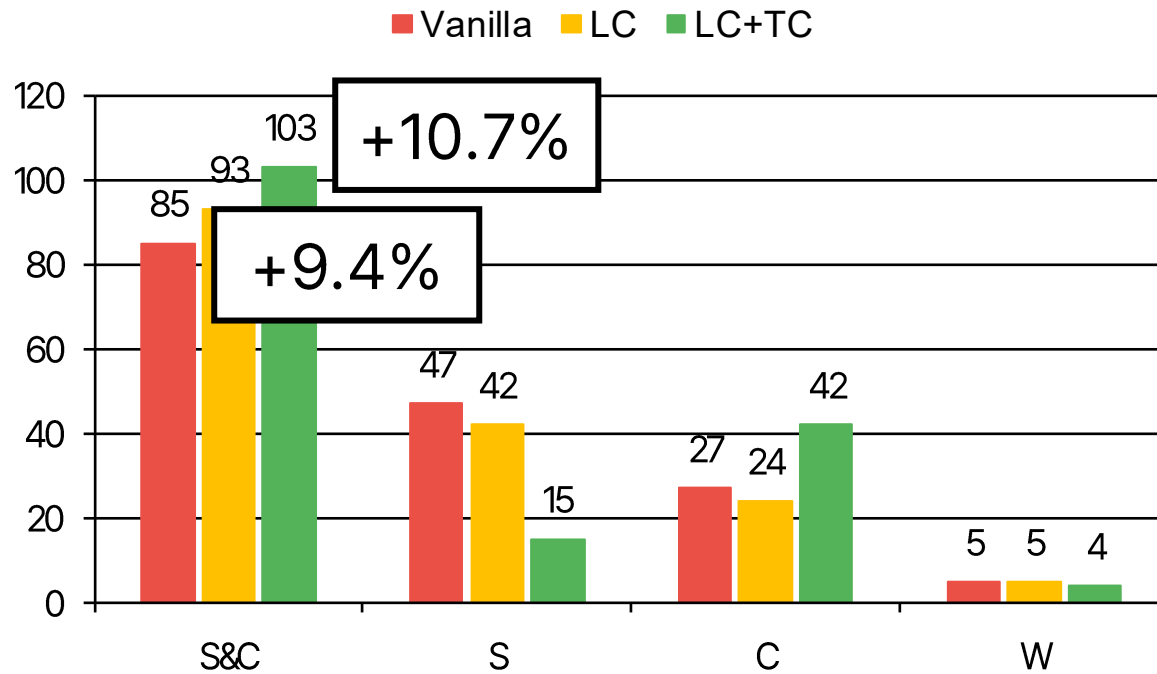
Vanilla LC LC+TC

Vanilla LC LC+TC

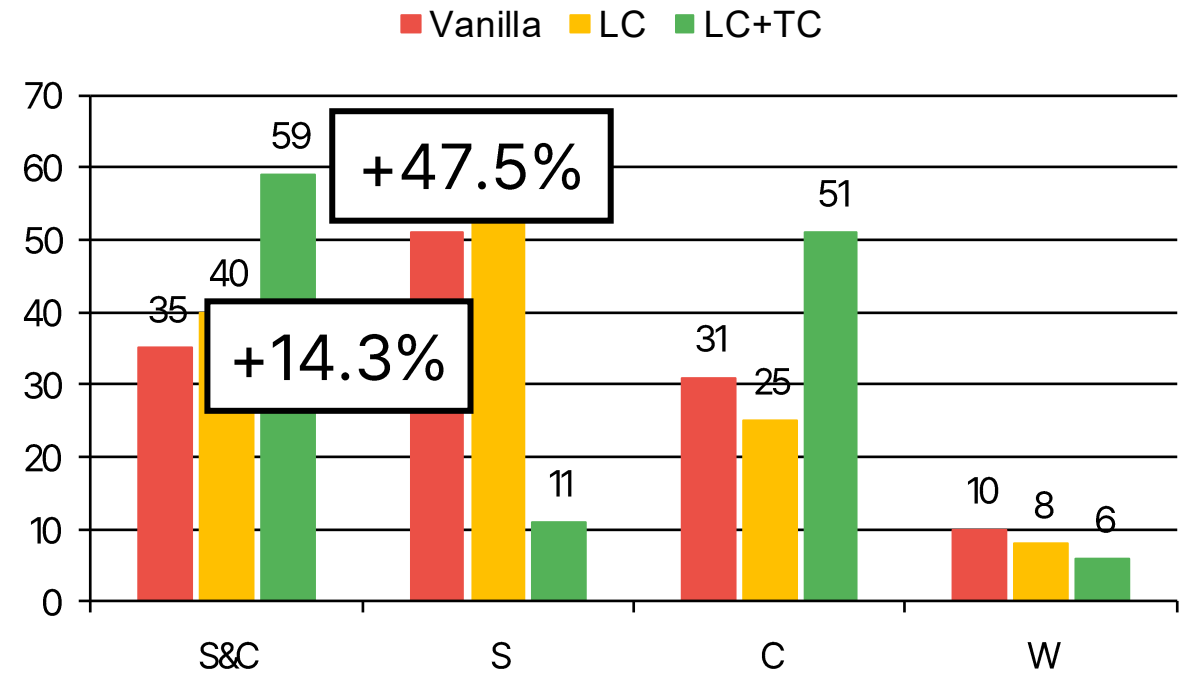


Impact of Checkers

HumanEval+



APPS



The More powerful checker derives the more S&C specs

```

def Spec(sq1: list[tuple[int, int]], sq2: list[tuple[int, int]], res: string) {
  var intersect =
     $\exists v :: (v \text{ in } sq1 \wedge \text{InPolygon}(sq2, v)) \vee \exists v :: (v \text{ in } sq2 \wedge \text{InPolygon}(sq1, v)) \vee \text{EdgesIntersect}(sq1, sq2)$ 
  if intersect then (res == "YES") else (res == "NO")
}

def InPolygon(poly: list[tuple[int, int]], pt: tuple[int, int]) {
  var n = len(poly);
   $\exists i :: (0 \leq i < n \wedge \text{OnSegment}(\text{poly}[i], \text{pt}, \text{poly}[(i + 1) \% n])) \vee (\text{RayIntersections}(\text{poly}, \text{pt}) \% 2 == 1)$ 
}

def OnSegment(p: tuple[int, int], q: tuple[int, int], r: tuple[int, int]) {
  var px, py = p; var qx, qy = q; var rx, ry = r;
   $(qx - px) * (ry - py) == (qy - py) * (rx - px) \wedge$ 
   $qx \geq (if\ px < rx\ then\ px\ else\ rx) \wedge qx \leq (if\ px > rx\ then\ px\ else\ rx) \wedge$ 
   $qy \geq (if\ py < ry\ then\ py\ else\ ry) \wedge qy \leq (if\ py > ry\ then\ py\ else\ ry)$ 
}

def RayIntersections(poly: list[tuple[int, int]], pt: tuple[int, int]) -> (count: int) {
  var n = len(poly);
  ensure (count == fold(lambda (acc: int, i: int) =
    if RayIntersectsSegment(poly[i], poly[(i + 1) \% n], pt) then (acc + 1) else acc, 0, [0 .. (n - 1)]));
}

def RayIntersectsSegment(p1: tuple[int, int], p2: tuple[int, int], q: tuple[int, int]) {
  var px, py = q; var x1, y1 = p1; var x2, y2 = p2;
  var y_min = (if y1 < y2 then y1 else y2); var y_max = (if y1 > y2 then y1 else y2);
  var x_min = (if x1 < x2 then x1 else x2); var x_max = (if x1 > x2 then x1 else x2);
  var segment_horizontal = (y1 == y2);
  var point_on_horizontal_segment = (segment_horizontal  $\wedge$  (py == y1))  $\wedge$  (px >= x_min)  $\wedge$  (px <= x_max);
  var intersects_right =
    ( $\neg$  segment_horizontal)  $\wedge$  (py >= y_min)  $\wedge$  (py <= y_max)  $\wedge$ 
     $((x1 + ((py - y1) * (x2 - x1)) / (y2 - y1))) \geq px$ );
  (point_on_horizontal_segment  $\vee$  intersects_right)
}

def EdgesIntersect(p1: list[tuple[int, int]], p2: list[tuple[int, int]]) {
  var n1 = len(p1); var n2 = len(p2);
   $\exists (i: \text{int}, j: \text{int}) :: ((0 \leq i < n1) \wedge (0 \leq j < n2) \wedge \text{EdgePairIntersect}(p1[i], p1[(i + 1) \% n1], p2[j], p2$ 
     $[(j + 1) \% n2]))$ 
}

def EdgePairIntersect(p1: tuple[int, int], p2: tuple[int, int], q1: tuple[int, int], q2: tuple[int, int]) {
  var o1 = Orientation(p1, p2, q1); var o2 = Orientation(p1, p2, q2);
  var o3 = Orientation(q1, q2, p1); var o4 = Orientation(q1, q2, p2);
   $(o1 \neq o2 \wedge o3 \neq o4) \vee$ 
   $(o1 == 0 \wedge \text{OnSegment}(p1, q1, p2)) \vee (o2 == 0 \wedge \text{OnSegment}(p1, q2, p2)) \vee$ 
   $(o3 == 0 \wedge \text{OnSegment}(q1, p1, q2)) \vee (o4 == 0 \wedge \text{OnSegment}(q1, p2, q2))$ 
}

def Orientation(p: tuple[int, int], q: tuple[int, int], r: tuple[int, int]) -> (o: int) {
  var px, py = p; var qx, qy = q; var rx, ry = r;
  var val =  $(qy - py) * (rx - qx) - (qx - px) * (ry - qy)$ ;
  ensure (val == 0 ==> o == 0);
  ensure (val > 0 ==> o == 1);
  ensure (val < 0 ==> o == 2);
}

```

Fig. 13. Full definition of the generated specification for the example problem in Fig. 7